

# Lawrence Livermore National Laboratory

## Towards an Abstraction-Friendly Programming Model for High Productivity and High Performance Computing



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# Programming and high level abstractions

- Best practice in software engineering:
  - High level abstractions → high productivity
    - expose interface, hide implementation details
    - + code reuse, - software complexity
  - Standard or customized
    - C++ STL containers, algorithms, iterators, ...
    - User-defined classes, functions, libraries, ...



# Semantics of abstractions: an unexploited gold mine

- Semantics: any standard or user-defined meanings
  - `STL::vector<T>` elements stored contiguously
  - `a->foo(x)` read only
  - `STL::Set<mytype>` order does not matter
- Not fully exploited by traditional programming models:
  - Traditional pmodel: write code and throw it to a vendor compiler
  - Low level intermediate representation (IR)
  - Info. hiding mechanism
  - Semantics: too many and too diverse
- Performance is inversely proportional to the level of abstractions used

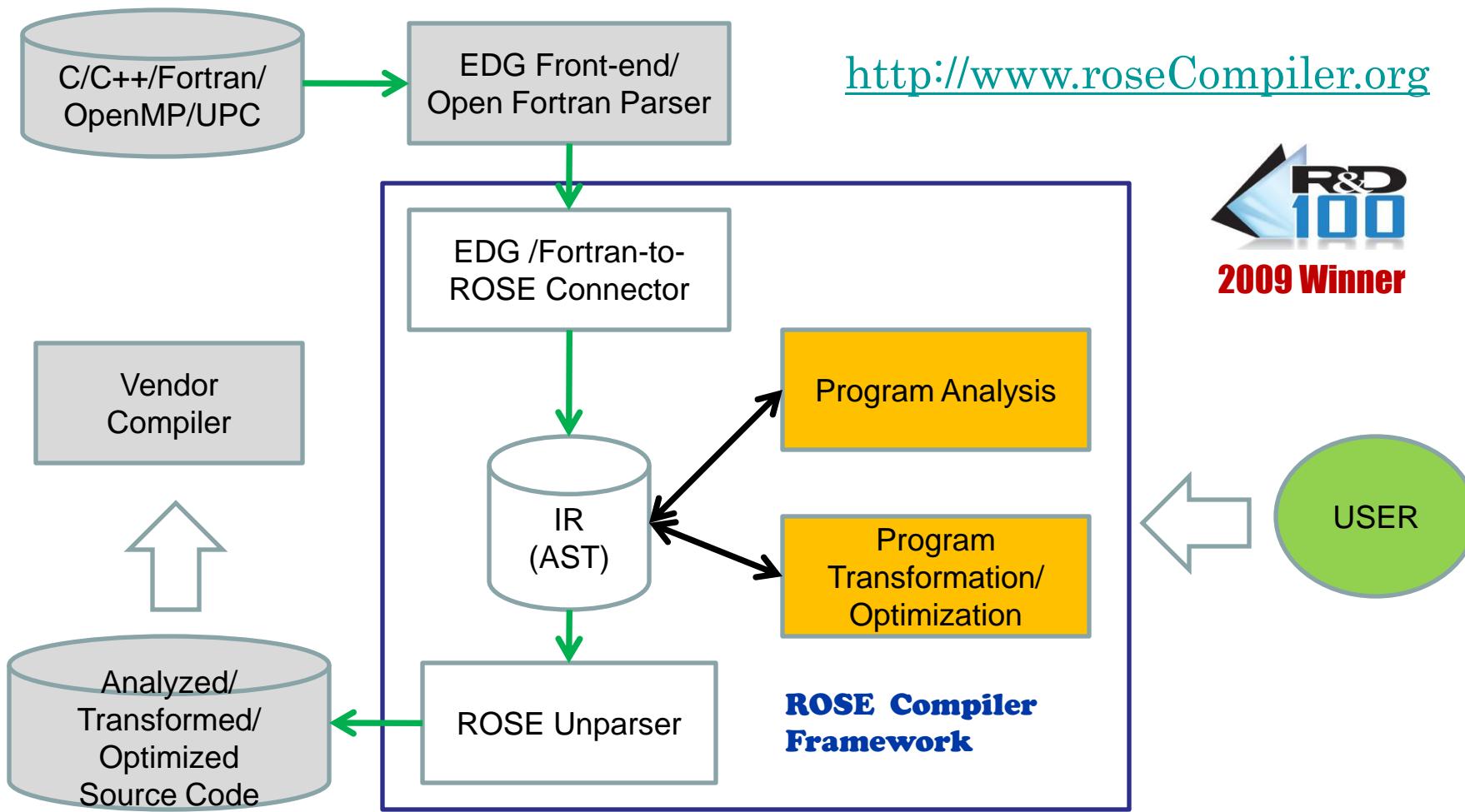


# An abstraction-friendly HPC programming model

- Goal: encourage the best programming practice while maintaining high or even better performance
- Solution:
  - User intervention
    - A specification: abstractions + semantics
    - User-defined optimizations: eliminate the dependence on compiler experts
  - An extensible source-to-source compiler framework
    - Recognize abstractions
    - Semantics-aware optimizations
    - Complement vendor compilers



# ROSE: making compiler technology accessible



# ROSE intermediate representation (IR)

- ROSE IR = AST + symbol tables + CFG ...
- Preserves all source level details
  - Token stream, source comments
  - C preprocessor control structure
  - C++ templates
- Rich interface
  - AST traversal, query, creation, copy, symbol lookup
  - Generic analyses, transformations, optimizations
- Fully support
  - Abstraction recognition and semantic analysis



# Case 1: a vector computation loop

```
std::vector <double> v1(SIZE);  
...  
double sum = 0.0;  
std::vector <double>::iterator iter;  
for (iter= v1.begin(); iter!=v1.end(); iter++)  
    sum = sum + *iter;
```

- std::vector : better productivity than arrays
  - Dynamic allocated, automatic de-allocation
  - Easier resizing, boundary-check
- Impede optimizations
  - E.g. Auto parallelization: primitive arrays
  - Non-canonical loop: *for (integer-init; test; increment) block*
  - Obscure element accesses: dereferencing an iterator



# Semantics can help

- Semantics of *std::vector <T>*
  - An array-like container
  - Element access methods: *[]*, *at()*, *\*iterator*
- Loop normalization:
  - loops using iterators → loops using integers
  - *\*iterator → container.at(i)*
- Dependence analysis:
  - *element\_access\_method(i) → subscript i*

```
double v1[SIZE];
double sum = 0.0;
int i;
#pragma omp parallel for reduction (sum)
for (i = 0; i<SIZE; i++)
    sum = sum + v1.at(i);
```



# Case 2: a domain-specific tree traversal

- Compass: A ROSE-based tool for static code analysis
  - A checker: detect a violation of MISRA\* Rule 5-0-18

```
void CompassAnalyses::PointerComparison::Traversal::visit(SgNode* node){  
    // Check binary operation nodes  
    SgBinaryOp* bin_op = isSgBinaryOp(node);  
    if (bin_op) {  
        // Check relational operations  
        if (isSgGreaterThanOp(node) || isSgGreaterOrEqualOp(node) ||  
            isSgLessThanOp(node) || isSgLessOrEqualOp(node)) {  
  
            SgType* lhs_type = bin_op->get_lhs_operand()->get_type();  
            SgType* rhs_type = bin_op->get_rhs_operand()->get_type();  
            // Check operands of pointer types  
            if (isSgPointerType(lhs_type) || isSgPointerType(rhs_type))  
                // output a violation  
                output->addOutput(bin_op);  
    } } }
```

\*The Motor Industry Software Reliability Association: MISRA C++: 2008 Guidelines for the use of the C++ language in critical systems.



# Semantics can help

- Enabling parallelization
  - Read-only semantics
    - Information retrieval functions: `get_*`()
    - Type casting functions: `isSg*`()
  - Order independent side effects
    - `output->addOutput(bin_op)`
    - Suitable for using `omp critical`
- Enabling customized optimization
  - Order-independent tree traversal
  - Nodes stored in memory pools
  - Recursive tree traversal → Loop over memory pools



# Optimized and Parallelized Code

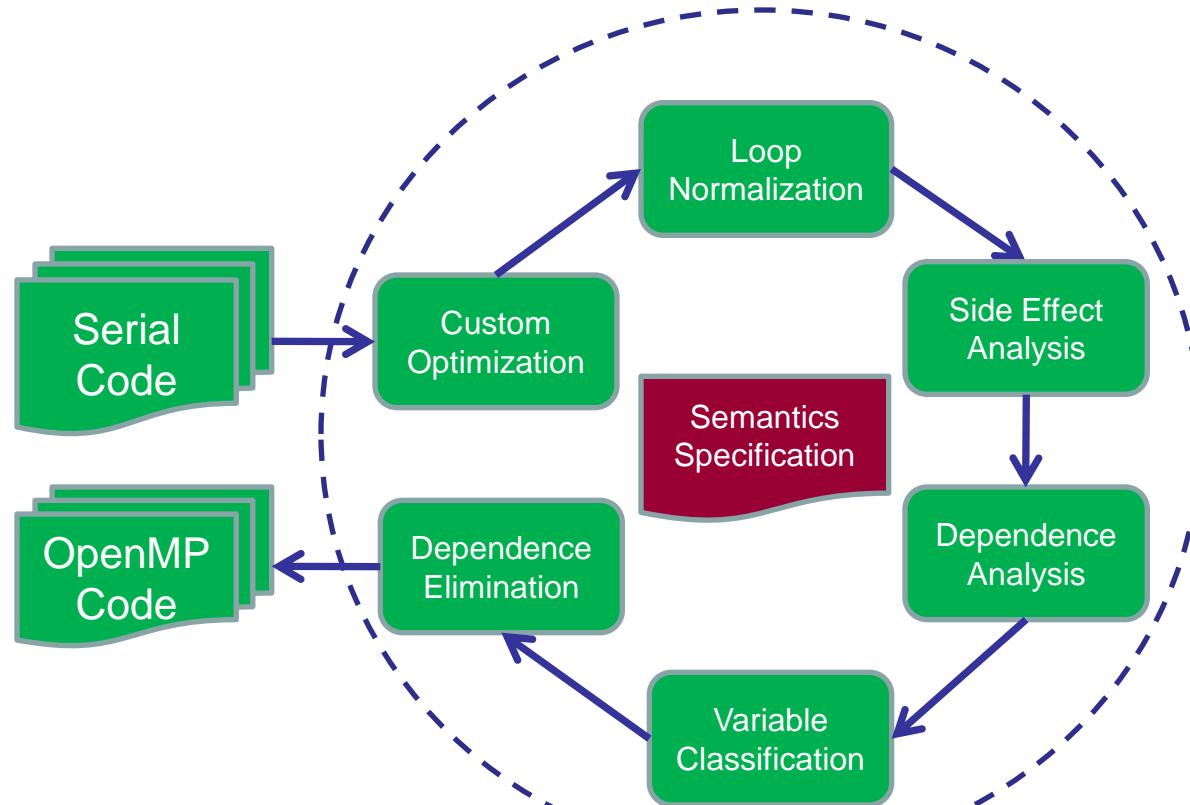
```
#pragma omp parallel for
for (i=0; i<pool_size; i++)
{
    SgBinaryOp* bin_op = isSgBinaryOp(MEMPOOL[i]);
    if (bin_op)
    {
        if (isSgGreaterThanOrEqualOp(node) || isSgGreaterOrEqualOp(node) ||
            isSgLessThanOp(node) || isSgLessOrEqualOp(node)) {

            SgType* lhs_type = bin_op->get_lhs_operand()->get_type();
            SgType* rhs_type = bin_op->get_rhs_operand()->get_type();

            if (isSgPointerType(lhs_type) || isSgPointerType(rhs_type))
            {
                #pragma omp critical
                output->addOutput(bin_op);
            }
        }
    }
}
```



# Implementation: a semantics-aware parallelizer



Chunhua Liao, Daniel J. Quinlan, Jeremiah J. Willcock and Thomas Panas, **Extending Automatic Parallelization to Optimize High-Level Abstractions for Multicore**, *IWOMP 2009 - International Workshop on OpenMP*, Dresden, Germany, 3-5 June 2009

# An abstraction/semantics specification file

```
class std::vector<MyType> {
    alias none; overlap none; //elements are alias-free and non-overlapping
    is_fixed_sized_array { //semantic-preserving functions as a fixed-sized array
        length(i) = {this.size()};
        element(i) = {this.operator[](i); this.at(i);};
    };
};

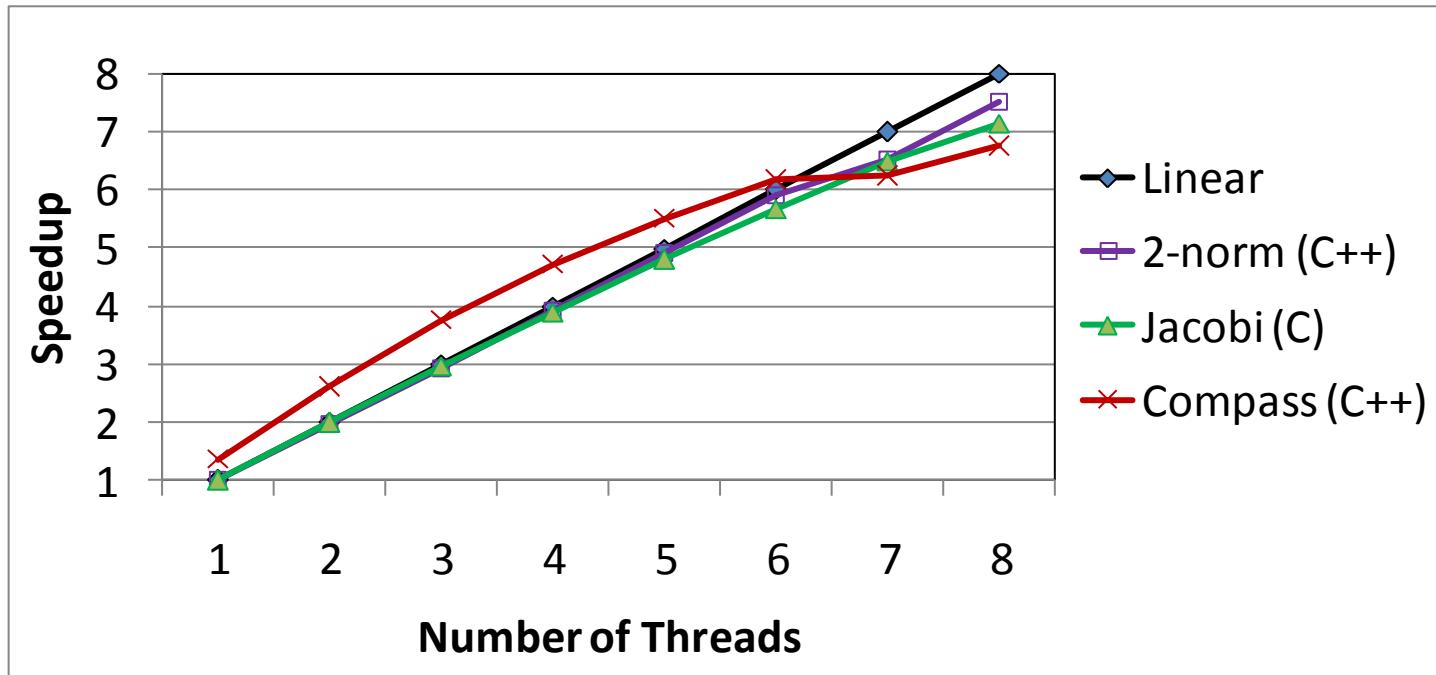
SgXXX* isSgXXX(SgNode*node)
{ modify none; } // read-only functions

SgNode* SgNode::get_XXX()
{ modify none; } // read-only member functions

void Compass::OutputObject::addOutput(SgNode* node){
    read {node};
    //order-independent side effects
    modify {Compass::OutputObject::outputList<order_independent>};
}
```



# Preliminary results



Platform: Dell Precision T5400, 3.16GHz quad-core Xeon X5460 dual processor, 8GB  
Compilers: ROSE OpenMP translator + Omni 1.6 Runtime + GCC 4.1.2

# Related Work

- Kennedy, et.al. Telescoping languages: a system for automatic generation of domain languages, proceedings of IEEE, 2005
  - High-level scripting languages, library preprocessing
- Gregor, Schupp, STLLint: lifting static checking from languages to libraries, Softw. Pract. Exper. 2006
  - Static analysis of error use of abstractions
  - C++ syntax for specification
- Kulkarni, Pingali, et.al. Optimistic parallelism requires abstractions. PLDI 2007
  - Abstraction: un-ordered set;
  - Semantics: commutativity, inverse



# Conclusions and future work

- ROSE-based abstraction-friendly programming model:
  - High productivity (use of abstractions) + high performance (semantics-aware optimizations)
  - Source-to-source: complement vendor compilers
  - User intervention: less depend on compiler experts
- Future work
  - Better specification files
  - Classify and formalize more abstractions/semantics
  - Operations on semantics to generate new semantics?



# Questions?

